### **Introduction to Local Search**

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### **Outline**

- Decision space
- Neighbourhood structure
- Neighbourhood exploration
  - Local move enumeration
  - Random sampling with replacement
  - Random sampling without replacement





The decision space is the set of all candidate solutions that may be visited during the solving process

- Only *feasible* solutions are considered in local search
  - ⇒ The objective function is defined on the whole decision space
  - The internal solution structure may not be specified
  - Solutions are *points* in decision space
- Initial feasible solution(s) may be:
  - Generated at random
  - Obtained by constructive search
  - Obtained heuristically in other ways





### **Examples**

How to generate solutions *uniformly at random* for the following problems?

• Travelling salesman





### **Examples**

- Travelling salesman
  - Generate a random permutation of the *N* cities
- Cable-trench problem (solutions are spanning trees)





#### **Examples**

- Travelling salesman
  - Generate a random permutation of the *N* cities
- Cable-trench problem (solutions are spanning trees)
  - By Cayley's theorem, there are  $N^{N-2}$  labeled spanning trees in a complete graph with N nodes
  - Generate a vector of N-2 integers in  $\{1,...,N\}$  uniformly at random
  - Decode the resulting *Prüfer sequence* into a spanning tree in linear time (see https://doi.org/10.4236/jsea.2009.22016)





### **Examples**

How to generate solutions *uniformly at random* for the following problems?

• Binary knapsack problem





### **Examples**

- Binary knapsack problem
  - Generate a random binary vector of length equal to the number of available items and reject sets of items that do not meet the capacity constraint(s)
  - <u>∧</u> Possibly (very) inefficient if the problem is tightly constrained
- Clique partitioning problem





#### **Examples**

- Binary knapsack problem
  - Generate a random binary vector of length equal to the number of available items and reject sets of items that do not meet the capacity constraint(s)
  - ↑ Possibly (very) inefficient if the problem is tightly constrained
- Clique partitioning problem
  - Draw a number of bins at random according to a *specific* distribution (see https://doi.org/10.1016/0097-3165(83)90009-2)
  - Randomly assign each vertex to one of those bins (some bins may end up empty)





# **Neighbourhood structure**

What are similar solutions?

- "Parts" of the two solutions are somehow identical
- Similar performance (in most cases)
- Connect the whole space

### **Example 1**

Symmetric Travelling Salesman Problem

- Tour length is the sum of the lengths of the tour edges
- Solutions are similar if they differ in a small number of edges
- 2-opt and 3-opt moves





## **Neighbourhood structure**

#### **Example 2**

Asymmetric Travelling Salesman Problem

- Tour length is the sum of the lengths of the tour arcs
- Solutions are similar if they differ in a small number of arcs
- Due to asymmetry, 2-opt moves and some 3-opt moves are not suitable because they reverse parts of the tour
- Insertion move as a special case of 3-opt

**Note:** What about neighbourhood size?





#### Local search in a nutshell

- 1. Visit neighbours of a current solution
- 2. Decide whether to reject them or to accept one as the next solution
- 3. Repeat

### Local move generation

- Enumeration
- Random sampling with replacement
- Random sampling without replacement





#### **Local move enumeration**

- Full neighbourhood exploration
- Enumeration order dictated by convenience
- Filter out invalid moves if needed
- Neighbourhood size may not be known in advance

### Random sampling with replacement

- Sampling uniformly at random typically preferred
- Rejection sampling useful when neighbourhood size not known in advance





#### Random sampling without replacement

- Partial to full neighbourhood exploration
- Random enumeration order to avoid search bias
- Filter out invalid moves if needed (rejection)
- Neighbourhood size may not be known in advance
- Generating (pseudo-)random permutations
  - Fisher-Yates shuffle
  - Linear congruential generator (low quality but constant space)
  - A plethora of other approaches





### Random sampling without replacement

```
Fisher-Yates shuffle
def fisher_yates_iter(n):
    p = list(range(n))
    for i in range(n-1, -1, -1):
        r = random.randrange(i+1)
        yield p[r]
        \# p[r], p[i] = p[i], p[r] \# preserve the permutation
        p[r] = p[i] # lazy
```





### Random sampling without replacement

Sparse Fisher-Yates shuffle

```
def sparse_fisher_yates_iter(n):
    p = dict()
    for i in range(n-1, -1, -1):
        r = random.randrange(i+1)
        yield p.get(r, r)
        if i != r:
            \# p[r] = p.pop(i, i) \# saves memory, takes time
            p[r] = p.get(i, i) # lazy, but faster
```





# **Concluding remarks**

- Feasible solutions are just points in decision space
- Internal solution structure is not exposed
- Neighbourhood structure should
  - Capture the similarity of solutions
  - Connect the whole decision space
- Neighbourhood enumeration and random sampling





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